

Зависимые типы в GHC 8

Максим Талдыкин

max@formalmethods.ru

Что такое?

Just 1 : Maybe Int : * : □

Maybe : * -> * : □

Vec Int Z : * : □

Vec : * -> Nat -> * : □

map : forall a b. (a -> b) -> [a] -> [b]
map f (x:xs) = f x : map f xs
map f [] = []

```
map :: forall a b. (a -> b) -> [a] -> [b]
map =
  \ (@a)
    (@b)
    (f :: a -> b)
    (ds :: [a]) ->
  case ds of
    [] -> GHC.Types.[] @b
  : x xs ->
    GHC.Types.: @b (f x) (map @a @b f xs)
```

forall

(a :: *)

(b :: *)

(f :: a -> b)

(xs :: [a])

. [b]

```
replicate :: pi (n : Nat) -> a -> Vec a n
replicate Z      _ = Nil
replicate (S m) x = x :> replicate m a
```


	vis	dep	rel
forall (a : *). a	-	+	-
a -> a	+	-	+
Monad m => m a	-	-	+
pi (a : Bool). f a	-	+	+
pi (a : Bool) -> f a	+	+	+

*** : ***

pi

Зачем?

zipWithN, mapT

tail []

strncpy

```
  :: forall (p q :: Nat)
=> (LE n p ~ True, LE n q ~ True)
  . pi (n :: Nat)
  -> Ptr Char p
  -> Ptr Char q
  -> IO ()
```

- [\$\Pi\$ -Ware: Hardware Description](#)
 - Swierstra et al., 2015
- [Correct-by-Construction Concurrency](#)
 - Brady & Hammond, 2009
- [Security-Typed Programming](#)
 - Morgenstern & Licata, 2010

Type Families

```
data Nat = Z | S Nat
```

```
type family
```

```
  (m :: Nat) :+ (n :: Nat) :: Nat
```

```
where
```

```
  Z      :+ n = n
```

```
  (S k) :+ n = S (k :+ n)
```

```
-- Vec a m -> Vec a n -> Vec a (m :+ n)
```


GADTs

Generalized Algebraic Data Types

```
data Vec :: Nat -> * -> *  
  Nil   :: Vec Z a  
  (::>) :: a -> Vec n a -> Vec (S n) a
```

```
data Vec (n :: Nat) (a :: *) :: * where
  Nil    :: n ~ Z => Nil n a
  (:>)  :: forall (m :: Nat)
    . n ~ S m
    => a -> Vec m a -> Vec n a
```

```
tail :: Vec (S k) a -> Vec k a
tail (_ :> xs) = xs
  -- (m :: Nat) (S k ~ S m)
```

type family

(v :: Vec a m) :++ (w :: Vec a n) :: Vec a k

where

Nil :++ w = w

(x :> xs) :++ w = x :> (xs :++ w)

type family

$(v :: \mathbf{Vec} \ a \ m) \ :++ \ (w :: \mathbf{Vec} \ a \ n) \ :: \mathbf{Vec} \ a \ k$

where

$\mathbf{Nil} \quad \quad \quad \ :++ \ w = w$

$(x \ :> \ xs) \ :++ \ w = x \ :> \ (xs \ :++ \ w)$

'Vec' of kind '* -> Nat -> *' is not promotable

In the kind **'Vec a m'**

Vec :: $\square \rightarrow \square \rightarrow \square$

```
replicate :: pi (n :: Nat) -> a -> Vec a n
replicate Z a = Nil
replicate (S m) a = a :> replicate m a
```



```
data Nat = Z | S Nat
```

```
data Nat's :: Nat -> * where  
  Z's :: Nat's Z  
  S's :: Nat's n -> Nat's (S n)
```

```
-- S's Z's :: Nat's (S Z)
```

```
replicate :: Nat's n -> a -> Vec a n
replicate Z's a = Nil
replicate (S's m) a = a :> replicate m a
```

Они уже здесь!

Singleton types here

Singleton types there

Singleton types everywhere

Monnier & Haguenaue, 2009

Hasochism

Lindley & mcBride, 2013

```
data Fin :: Nat -> * where
  FZ :: Fin (S n)
  FS :: Fin n -> Fin (S n)
```

```
lookup :: pi (f :: Fin n) -> Vec a n -> a
lookup FZ (x :> _) = x
lookup (FS i) (_ :> xs) = lookup i xs
```

```
data Fin's (n :: Nat) (f :: Fin n) :: * where
  FZ's :: Fin's (S m) FZ
  FS's :: Fin's m g -> Fin's (S m) (FS g)

-- FS's FZ's :: Fin's 4 (FS FZ :: Fin 4)
```



```
data Fin's (n :: Nat) (f :: Fin n) :: * where  
  FZ's :: Fin's (S m) FZ  
  FS's :: Fin's m g -> Fin's (S m) (FS g)
```

Kind variable also used **as type** variable: 'n'
In the **data type** declaration for '**Fin's**'

```
get "/Contract/:id" $ do
  intParam "id" >>= queryDb "Contract" >>= json
```

```
type API
  = "Contract"
  :> Capture "id" Int
  :> Get '[JSON] Contract
```

```
type Api
  = "Contract"
  :> Capture "id" Int
  :> RoleFilter "Contract" "owner" '[43, 265]
  :> Get '[JSON] (Obj Contract)
```

```
data User (fieldName :: Symbol) where
  Id :: Int -> User "id"
  Name :: Text -> User "name"
  Roles :: [Vector (Ref Role)] -> User "roles"
```

```
data Contract (fieldName :: Symbol) where
  Id :: Int -> Contract "id"
  Owner :: Ref Role -> Contract "owner"
```

```
type family TableName (m :: Symbol -> *) :: Symbol
type instance TableName User = "User"
type instance TableName Contract = "Contract"
```

```
type Api
  = "Contract"
    :> Capture "id" Int
    :> RoleFilter "Contract" "owner" '[43, 265]
    :> Get '[JSON] (Obj Contract)
```

```
type Api
  = MkFilter Contract
    (RoleFilter C.Owner '[Role1, Role2])
```

```
type family MkFilter  
  (f :: Symbol)  
  (m :: Symbol -> *)  
  (filter :: *)  
where  
  MkFilter f m (RoleFilter (own :: m g) rs)  
    = TableName m  
      :> Capture  
        (FieldName (TableId m))  
        (FieldType (TableId m))  
      :> RoleFilter' (TableName m) g rs  
      :> Get '[JSON] (Obj m)
```

Termination

Equality

Consistency

- [System FC with Explicit Kind Equality](#), 2013
- [Dependent Types in Haskell](#), draft
- [Type Inference, Haskell and Dependent Types](#), 2013



Stephanie Weirich



Richard A. Eisenberg



Per Martin-Löf, [Intuitionistic type theory \(1984\)](#)

FORM \forall L M \exists T H O D S

